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OWL 2 Web Ontology Language
Direct Semantics

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Abstract

The OWL 2 Web Ontology Language, informally OWL 2, is an ontology language for the Semantic Web with formally defined meaning. OWL 2 ontologies provide classes, properties, individuals, and data values and are stored as Semantic Web documents. OWL 2 ontologies can be used along with information written in RDF, and OWL 2 ontologies themselves are primarily exchanged as RDF documents. The OWL 2 Document Overview describes the overall state of OWL 2, and should be read before other OWL 2 documents.
This document provides the direct model-theoretic semantics for OWL 2, which is compatible with the description logic *SROIQ*. Furthermore, this document defines the most common inference problems for OWL 2.

Status of this Document

May Be Superseded

This section describes the status of this document at the time of its publication. Other documents may supersede this document. A list of current W3C publications and the latest revision of this technical report can be found in the W3C technical reports index at http://www.w3.org/TR/.

Summary of Changes

This Last Call Working Draft has a few changes since the previous version of 02 December 2008.

- Several changes reflect changes to surface structure of the functional syntax.
- Several changes reflect changes to the treatment of datatypes.

(Second) Last Call

The Working Group believes it has completed its design work for the technologies specified this document, so this is a "Last Call" draft. The design is not expected to change significantly, going forward, and now is the key time for external review, before the implementation phase. (This is the second Last Call draft of this document. The public response to the previous Last Call prompted the Working Group to make material changes to the design.)

Please Comment By 12 May 2009

The OWL Working Group seeks public feedback on this Working Draft. Please send your comments to public-owl-comments@w3.org (public archive). If possible, please offer specific changes to the text that would address your concern. You may also wish to check the Wiki Version of this document and see if the relevant text has already been updated.

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1 Introduction

This document defines the direct model-theoretic semantics of OWL 2. The semantics given here is strongly related to the semantics of description logics [*Description Logics*] and it extends the semantics of the description logic *SROIQ* [*SROIQ*]. As the definition of *SROIQ* does not provide for datatypes and punning, the semantics of OWL 2 is defined directly on the constructs of the structural specification of OWL 2 [[OWL 2 Specification](#)] instead of by reference to *SROIQ*. For the constructs available in *SROIQ*, the semantics of *SROIQ* trivially corresponds to the one defined in this document.

Since each OWL 1 DL ontology is an OWL 2 ontology, this document also provides a direct semantics for OWL 1 Lite and OWL 1 DL ontologies; this semantics is equivalent to the direct model-theoretic semantics of OWL 1 Lite and OWL 1 DL [[OWL Abstract Syntax and Semantics](#)]. Furthermore, this document also provides the direct model-theoretic semantics for the OWL 2 profiles [[OWL 2 Profiles](#)].

The semantics is defined for OWL 2 axioms and ontologies, which should be understood as instances of the structural specification [[OWL 2 Specification](#)]. Parts of the structural specification are written in this document using the functional-style syntax.

OWL 2 allows ontologies, anonymous individuals, and axioms to be annotated; furthermore, annotations themselves can contain additional annotations. All these types of annotations, however, have no semantic meaning in OWL 2 and are ignored in this document. OWL 2 declarations are used only to disambiguate class expressions from data ranges and object property from data property expressions in the functional-style syntax; therefore, they are not mentioned explicitly in this document.

2 Direct Model-Theoretic Semantics for OWL 2

This section specifies the direct model-theoretic semantics of OWL 2 ontologies.

2.1 Vocabulary

A *datatype map*, formalizing [datatype maps](#) from the OWL 2 Specification [[OWL 2 Specification](#)], is a 6-tuple $D = (N_{DT} , N_{LS} , N_{FS} , \cdot^{DT} , \cdot^{LS} , \cdot^{FS})$ with the following components:

- N_{DT} is a set of datatypes (more precisely, names of datatypes) that does not contain the datatype *rdfs:Literal*.
- N_{LS} is a function that assigns to each datatype $DT \in N_{DT}$ a set $N_{LS}(DT)$ of strings called *lexical forms*. The set $N_{LS}(DT)$ is called the *lexical space* of *DT*.
- N_{FS} is a function that assigns to each datatype $DT \in N_{DT}$ a set $N_{FS}(DT)$ of pairs $\langle F , v \rangle$, where *F* is a *constraining facet* and *v* is an arbitrary data value called the *constraining value*. The set $N_{FS}(DT)$ is called the *facet space* of *DT*.
- For each datatype $DT \in N_{DT}$, the *interpretation function* $\cdot^{DT}$ assigns to *DT* a set $(DT)^{DT}$ called the *value space* of *DT*.
- For each datatype $DT \in N_{DT}$ and each lexical form $LV \in N_{LS}(DT)$, the *interpretation function* $\cdot^{LS}$ assigns to the pair $\langle LV , DT \rangle$ a *data value* $(\langle LV , DT \rangle)^{LS} \in (DT)^{DT}$.

- For each datatype $DT \in N_{DT}$ and each pair $\langle F , v \rangle \in N_{FS}(DT)$, the *interpretation function* $\cdot^{FS}$ assigns to $\langle F , v \rangle$ the set $(\langle F , v \rangle)^{FS} \subseteq (DT)^{DT}$.

A *vocabulary* $V = (V_C , V_{OP} , V_{DP} , V_I , V_{DT} , V_{LT} , V_{FA})$ over a datatype map D is a 7-tuple consisting of the following elements:

- V_C is a set of *classes* as defined in the OWL 2 Specification [*OWL 2 Specification*], containing at least the classes *owl:Thing* and *owl:Nothing*.
- V_{OP} is a set of *object properties* as defined in the OWL 2 Specification [*OWL 2 Specification*], containing at least the object properties *owl:topObjectProperty* and *owl:bottomObjectProperty*.
- V_{DP} is a set of *data properties* as defined in the OWL 2 Specification [*OWL 2 Specification*], containing at least the data properties *owl:topDataProperty* and *owl:bottomDataProperty*.
- V_I is a set of *individuals* (named and anonymous) as defined in the OWL 2 Specification [*OWL 2 Specification*].
- V_{DT} is a set containing all datatypes of D, the datatype *rdfs:Literal*, and possibly other datatypes; that is, $N_{DT} \cup \{ rdfs:Literal \} \subseteq V_{DT}$.
- V_{LT} is a set of *literals* LV^^DT for each datatype $DT \in N_{DT}$ and each lexical form $LV \in N_{LS}(DT)$.
- V_{FA} is the set of pairs $\langle F , lt \rangle$ for each constraining facet F, datatype $DT \in N_{DT}$, and literal $lt \in V_{LT}$ such that $\langle F , (\langle LV , DT_1 \rangle)^{LS} \rangle \in N_{FS}(DT)$, where LV is the lexical form of lt and DT_1 is the datatype of lt.

Given a vocabulary V, the following conventions are used in this document to denote different syntactic parts of OWL 2 ontologies:

- `OP` denotes an object property;
- `OPE` denotes an object property expression;
- `DP` denotes a data property;
- `DPE` denotes a data property expression;
- `C` denotes a class;
- `CE` denotes a class expression;
- `DT` denotes a datatype;
- `DR` denotes a data range;
- `a` denotes an individual (named or anonymous);
- `lt` denotes a literal; and
- `F` denotes a constraining facet.

2.2 Interpretations

Given a datatype map D and a vocabulary V over D, an *interpretation* $I = (\Delta_I , \Delta_D , \cdot^C , \cdot^{OP} , \cdot^{DP} , \cdot^I , \cdot^{DT} , \cdot^{LT} , \cdot^{FA})$ for D and V is a 9-tuple with the following structure:

- Δ_I is a nonempty set called the *object domain*.
- Δ_D is a nonempty set disjoint with Δ_I called the *data domain* such that $(DT)^{DT} \subseteq \Delta_D$ for each datatype $DT \in V_{DT}$.

- $\cdot^{C}$ is the *class interpretation function* that assigns to each class $C \in V_C$ a subset $(C)^C \subseteq \Delta_I$ such that
  - $(owl{:}Thing)^C = \Delta_I$ and
  - $(owl{:}Nothing)^C = \emptyset$.
- $\cdot^{OP}$ is the *object property interpretation function* that assigns to each object property $OP \in V_{OP}$ a subset $(OP)^{OP} \subseteq \Delta_I \times \Delta_I$ such that
  - $(owl{:}topObjectProperty)^{OP} = \Delta_I \times \Delta_I$ and
  - $(owl{:}bottomObjectProperty)^{OP} = \emptyset$.
- $\cdot^{DP}$ is the *data property interpretation function* that assigns to each data property $DP \in V_{DP}$ a subset $(DP)^{DP} \subseteq \Delta_I \times \Delta_D$ such that
  - $(owl{:}topDataProperty)^{DP} = \Delta_I \times \Delta_D$ and
  - $(owl{:}bottomDataProperty)^{DP} = \emptyset$.
- $\cdot^{I}$ is the *individual interpretation function* that assigns to each individual $a \in V_I$ an element $(a)^I \in \Delta_I$.
- $\cdot^{DT}$ is the *datatype interpretation function* that assigns to each datatype $DT \in V_{DT}$ a subset $(DT)^{DT} \subseteq \Delta_D$ such that
  - $\cdot^{DT}$ is the same as in D for each datatype $DT \in N_{DT}$, and
  - $(rdfs{:}Literal)^{DT} = \Delta_D$.
- $\cdot^{LT}$ is the *literal interpretation function* that is defined as $(lt)^{LT} = (\langle LV , DT \rangle)^{LS}$ for each $lt \in V_{LT}$, where LV is the lexical form of lt and DT is the datatype of lt.
- $\cdot^{FA}$ is the *facet interpretation function* that is defined as $(\langle F , lt \rangle)^{FA} = (\langle F , (lt)^{LT} \rangle)^{FS}$ for each $\langle F , lt \rangle \in V_{FA}$.

The following sections define the extensions of $\cdot^{OP}$, $\cdot^{DT}$, and $\cdot^{C}$ to object property expressions, data ranges, and class expressions.

2.2.1 Object Property Expressions

The object property interpretation function $\cdot^{OP}$ is extended to object property expressions as shown in Table 1.

Table 1. Interpreting Object Property Expressions

Object Property Expression	Interpretation $\cdot^{OP}$
`ObjectInverseOf( OP )`	$\{ \langle x , y \rangle \mid \langle y , x \rangle \in (OP)^{OP} \}$

2.2.2 Data Ranges

The datatype interpretation function $\cdot^{DT}$ is extended to data ranges as shown in Table 3. All datatypes in OWL 2 are unary, so each datatype DT is interpreted as a unary relation over Δ_D — that is, as a set $(DT)^{DT} \subseteq \Delta_D$. OWL 2 currently does not define data ranges of arity more than one; however, by allowing for *n*-ary data ranges, the syntax of OWL 2 provides a "hook" allowing implementations to introduce extensions such as comparisons and arithmetic. An *n*-ary data range DR

is interpreted as an *n*-ary relation $(DR)^{DT}$ over Δ_D — that is, as a set $(DT)^{DT} \subseteq (\Delta_D)^n$

Table 3. Interpreting Data Ranges

Data Range	Interpretation $\cdot^{DT}$
`DataIntersectionOf( DR1 ... DRn )`	$(DR_1)^{DT} \cap \ldots \cap (DR_n)^{DT}$
`DataUnionOf( DR1 ... DRn )`	$(DR_1)^{DT} \cup \ldots \cup (DR_n)^{DT}$
`DataComplementOf( DR )`	$(\Delta_D)^n \setminus (DR)^{DT}$ where *n* is the arity of *DR*
`DataOneOf( lt1 ... ltn )`	$\{ (lt_1)^{LT} , \ldots , (lt_n)^{LT} \}$
`DatatypeRestriction( DT F1 lt1 ... Fn ltn )`	$(DT)^{DT} \cap (\langle F_1 , lt_1 \rangle)^{FA} \cap \ldots \cap (\langle F_n , lt_n \rangle)^{FA}$

2.2.3 Class Expressions

The class interpretation function $\cdot^C$ is extended to class expressions as shown in Table 4. For *S* a set, *#S* denotes the number of elements in *S*.

Table 4. Interpreting Class Expressions

Class Expression	Interpretation $\cdot^C$
`ObjectIntersectionOf( CE1 ... CEn )`	$(CE_1)^C \cap \ldots \cap (CE_n)^C$
`ObjectUnionOf( CE1 ... CEn )`	$(CE_1)^C \cup \ldots \cup (CE_n)^C$
`ObjectComplementOf( CE )`	$\Delta_I \setminus (CE)^C$
`ObjectOneOf( a1 ... an )`	$\{ (a_1)^I , \ldots , (a_n)^I \}$
`ObjectSomeValuesFrom( OPE CE )`	$\{ x \mid \exists y : \langle x, y \rangle \in (OPE)^{OP}$ and $y \in (CE)^C \}$
`ObjectAllValuesFrom( OPE CE )`	$\{ x \mid \forall y : \langle x, y \rangle \in (OPE)^{OP}$ implies $y \in (CE)^C \}$
`ObjectHasValue( OPE a )`	$\{ x \mid \langle x , (a)^I \rangle \in (OPE)^{OP} \}$
`ObjectHasSelf( OPE )`	$\{ x \mid \langle x , x \rangle \in (OPE)^{OP} \}$

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ObjectMinCardinality(n OPE)	{ x | #{ y | ⟨ x , y ⟩ ∈ $(OPE)^{OP}$ } ≥ n }
ObjectMaxCardinality(n OPE)	{ x | #{ y | ⟨ x , y ⟩ ∈ $(OPE)^{OP}$ } ≤ n }
ObjectExactCardinality(n OPE)	{ x | #{ y | ⟨ x , y ⟩ ∈ $(OPE)^{OP}$ } = n }
ObjectMinCardinality(n OPE CE)	{ x | #{ y | ⟨ x , y ⟩ ∈ $(OPE)^{OP}$ and y ∈ $(CE)^{C}$ } ≥ n }
ObjectMaxCardinality(n OPE CE)	{ x | #{ y | ⟨ x , y ⟩ ∈ $(OPE)^{OP}$ and y ∈ $(CE)^{C}$ } ≤ n }
ObjectExactCardinality(n OPE CE)	{ x | #{ y | ⟨ x , y ⟩ ∈ $(OPE)^{OP}$ and y ∈ $(CE)^{C}$ } = n }
DataSomeValuesFrom(DPE_1 ... DPE_n DR)	{ x | ∃ $y_1, \ldots , y_n$: ⟨ x , y_k ⟩ ∈ $(DPE_k)^{DP}$ for each 1 ≤ k ≤ n and ⟨ y_1 , ... , y_n ⟩ ∈ $(DR)^{DT}$ }
DataAllValuesFrom(DPE_1 ... DPE_n DR)	{ x | ∀ $y_1, \ldots , y_n$: ⟨ x , y_k ⟩ ∈ $(DPE_k)^{DP}$ for each 1 ≤ k ≤ n imply ⟨ y_1 , ... , y_n ⟩ ∈ $(DR)^{DT}$ }
DataHasValue(DPE lt)	{ x | ⟨ x , $(lt)^{LT}$ ⟩ ∈ $(DPE)^{DP}$ }
DataMinCardinality(n DPE)	{ x | #{ y | ⟨ x , y ⟩ ∈ $(DPE)^{DP}$} ≥ n }
DataMaxCardinality(n DPE)	{ x | #{ y | ⟨ x , y ⟩ ∈ $(DPE)^{DP}$ } ≤ n }
DataExactCardinality(n DPE)	{ x | #{ y | ⟨ x , y ⟩ ∈ $(DPE)^{DP}$ } = n }
DataMinCardinality(n DPE DR)	{ x | #{ y | ⟨ x , y ⟩ ∈ $(DPE)^{DP}$ and y ∈ $(DR)^{DT}$ } ≥ n }
DataMaxCardinality(n DPE DR)	{ x | #{ y | ⟨ x , y ⟩ ∈ $(DPE)^{DP}$ and y ∈ $(DR)^{DT}$ } ≤ n }
DataExactCardinality(n DPE DR)	{ x | #{ y | ⟨ x , y ⟩ ∈ $(DPE)^{DP}$ and y ∈ $(DR)^{DT}$ } = n }

2.3 Satisfaction in an Interpretation

An interpretation $I = (\Delta_I , \Delta_D , \cdot^{C} , \cdot^{OP} , \cdot^{DP} , \cdot^{I} , \cdot^{DT} , \cdot^{LT} , \cdot^{FA})$ *satisfies* an axiom w.r.t. an ontology O if the axiom satisfies the relevant condition from the following sections. Satisfaction of axioms in *I* is defined w.r.t. *O* because satisfaction of key axioms uses the following function:

> $ISNAMED_O(x) = true$ for $x \in \Delta_I$ if and only if $(a)^I = x$ for some named individual *a* occurring in the [axiom closure]() of *O*

2.3.1 Class Expression Axioms

Satisfaction of OWL 2 class expression axioms in *I* w.r.t. *O* is defined as shown in Table 5.

Table 5. Satisfaction of Class Expression Axioms in an Interpretation

Axiom	Condition
`SubClassOf( CE₁ CE₂ )`	$(CE_1)^C \subseteq (CE_2)^C$
`EquivalentClasses( CE₁ ... CEₙ )`	$(CE_j)^C = (CE_k)^C$ for each $1 \leq j \leq n$ and each $1 \leq k \leq n$
`DisjointClasses( CE₁ ... CEₙ )`	$(CE_j)^C \cap (CE_k)^C = \emptyset$ for each $1 \leq j \leq n$ and each $1 \leq k \leq n$ such that $j \neq k$
`DisjointUnion( C CE₁ ... CEₙ )`	$(C)^C = (CE_1)^C \cup \ldots \cup (CE_n)^C$ and $(CE_j)^C \cap (CE_k)^C = \emptyset$ for each $1 \leq j \leq n$ and each $1 \leq k \leq n$ such that $j \neq k$

2.3.2 Object Property Expression Axioms

Satisfaction of OWL 2 object property expression axioms in *I* w.r.t. *O* is defined as shown in Table 6.

Table 6. Satisfaction of Object Property Expression Axioms in an Interpretation

Axiom	Condition
`SubObjectPropertyOf( OPE₁ OPE₂ )`	$(OPE_1)^{OP} \subseteq (OPE_2)^{OP}$
`SubObjectPropertyOf( ObjectPropertyChain( OPE₁ ... OPEₙ ) OPE )`	$\forall y_0, \ldots, y_n : \langle y_0, y_1 \rangle \in (OPE_1)^{OP}$ and ... and $\langle y_{n-1}, y_n \rangle \in (OPE_n)^{OP}$ imply $\langle y_0, y_n \rangle \in (OPE)^{OP}$
`EquivalentObjectProperties( OPE₁ ... OPEₙ )`	$(OPE_j)^{OP} = (OPE_k)^{OP}$ for each $1 \leq j \leq n$ and each $1 \leq k \leq n$

`DisjointObjectProperties( OPE`$_1$ `... OPE`$_n$ `)`	$(OPE_j)^{OP} \cap (OPE_k)^{OP} = \emptyset$ for each $1 \leq j \leq n$ and each $1 \leq k \leq n$ such that $j \neq k$
`ObjectPropertyDomain( OPE CE )`	$\forall x, y : \langle x, y \rangle \in (OPE)^{OP}$ implies $x \in (CE)^C$
`ObjectPropertyRange( OPE CE )`	$\forall x, y : \langle x, y \rangle \in (OPE)^{OP}$ implies $y \in (CE)^C$
`InverseObjectProperties( OPE`$_1$ `OPE`$_2$ `)`	$(OPE_1)^{OP} = \{ \langle x, y \rangle \mid \langle y, x \rangle \in (OPE_2)^{OP} \}$
`FunctionalObjectProperty( OPE )`	$\forall x, y_1, y_2 : \langle x, y_1 \rangle \in (OPE)^{OP}$ and $\langle x, y_2 \rangle \in (OPE)^{OP}$ imply $y_1 = y_2$
`InverseFunctionalObjectProperty( OPE )`	$\forall x_1, x_2, y : \langle x_1, y \rangle \in (OPE)^{OP}$ and $\langle x_2, y \rangle \in (OPE)^{OP}$ imply $x_1 = x_2$
`ReflexiveObjectProperty( OPE )`	$\forall x : x \in \Delta_I$ implies $\langle x, x \rangle \in (OPE)^{OP}$
`IrreflexiveObjectProperty( OPE )`	$\forall x : x \in \Delta_I$ implies $\langle x, x \rangle \notin (OPE)^{OP}$
`SymmetricObjectProperty( OPE )`	$\forall x, y : \langle x, y \rangle \in (OPE)^{OP}$ implies $\langle y, x \rangle \in (OPE)^{OP}$
`AsymmetricObjectProperty( OPE )`	$\forall x, y : \langle x, y \rangle \in (OPE)^{OP}$ implies $\langle y, x \rangle \notin (OPE)^{OP}$
`TransitiveObjectProperty( OPE )`	$\forall x, y, z : \langle x, y \rangle \in (OPE)^{OP}$ and $\langle y, z \rangle \in (OPE)^{OP}$ imply $\langle x, z \rangle \in (OPE)^{OP}$

2.3.3 Data Property Expression Axioms

Satisfaction of OWL 2 data property expression axioms in *I* w.r.t. *O* is defined as shown in Table 7.

Table 7. Satisfaction of Data Property Expression Axioms in an Interpretation

Axiom	Condition
`SubDataPropertyOf( DPE`$_1$ `DPE`$_2$ `)`	$(DPE_1)^{DP} \subseteq (DPE_2)^{DP}$

`EquivalentDataProperties( DPE1 ... DPEn )`	$(DPE_j)^{DP} = (DPE_k)^{DP}$ for each $1 \leq j \leq n$ and each $1 \leq k \leq n$
`DisjointDataProperties( DPE1 ... DPEn )`	$(DPE_j)^{DP} \cap (DPE_k)^{DP} = \emptyset$ for each $1 \leq j \leq n$ and each $1 \leq k \leq n$ such that $j \neq k$
`DataPropertyDomain( DPE CE )`	$\forall x , y : \langle x , y \rangle \in (DPE)^{DP}$ implies $x \in (CE)^C$
`DataPropertyRange( DPE DR )`	$\forall x , y : \langle x , y \rangle \in (DPE)^{DP}$ implies $y \in (DR)^{DT}$
`FunctionalDataProperty( DPE )`	$\forall x , y_1 , y_2 : \langle x , y_1 \rangle \in (DPE)^{DP}$ and $\langle x , y_2 \rangle \in (DPE)^{DP}$ imply $y_1 = y_2$

2.3.4 Datatype Definitions

Satisfaction of datatype definitions in *I* w.r.t. *O* is defined as shown in Table 8.

Table 9. Satisfaction of Datatype Definitions in an Interpretation

Axiom	Condition
`DatatypeDefinition( DT DR )`	$(DT)^{DT} = (DR)^{DT}$

2.3.5 Keys

Satisfaction of keys in *I* w.r.t. *O* is defined as shown in Table 9.

Table 9. Satisfaction of Keys in an Interpretation

Axiom	Condition
`HasKey( CE ( OPE1 ... OPEm ) ( DPE1 ... DPEn ) )`	$\forall x , y , z_1 , \ldots , z_m , w_1 , \ldots , w_n$: if $x \in (CE)^C$ and $ISNAMED_O(x)$ and $y \in (CE)^C$ and $ISNAMED_O(y)$ and $\langle x , z_i \rangle \in (OPE_i)^{OP}$ and $\langle y , z_i \rangle \in (OPE_i)^{OP}$ and $ISNAMED_O(z_i)$ for each $1 \leq i \leq m$ and $\langle x , w_j \rangle \in (DPE_j)^{DP}$ and $\langle y , w_j \rangle \in (DPE_j)^{DP}$ for each $1 \leq j \leq n$ then x = y

2.3.6 Assertions

Satisfaction of OWL 2 assertions in *I* w.r.t. *O* is defined as shown in Table 10.

Table 10. Satisfaction of Assertions in an Interpretation

Axiom	Condition
`SameIndividual( a1 ... an )`	$(a_j)^I = (a_k)^I$ for each $1 \leq j \leq n$ and each $1 \leq k \leq n$
`DifferentIndividuals( a1 ... an )`	$(a_j)^I \neq (a_k)^I$ for each $1 \leq j \leq n$ and each $1 \leq k \leq n$ such that $j \neq k$
`ClassAssertion( CE a )`	$(a)^I \in (CE)^C$
`ObjectPropertyAssertion( OPE a1 a2 )`	$\langle (a_1)^I , (a_2)^I \rangle \in (OPE)^{OP}$
`NegativeObjectPropertyAssertion( OPE a1 a2 )`	$\langle (a_1)^I , (a_2)^I \rangle \notin (OPE)^{OP}$
`DataPropertyAssertion( DPE a lt )`	$\langle (a)^I , (lt)^{LT} \rangle \in (DPE)^{DP}$
`NegativeDataPropertyAssertion( DPE a lt )`	$\langle (a)^I , (lt)^{LT} \rangle \notin (DPE)^{DP}$

2.3.7 Ontologies

An interpretation *I* *satisfies* an OWL 2 ontology *O* if all axioms in the axiom closure of *O* (with anonymous individuals standardized apart as described in Section 5.6.2 of the OWL 2 Specification [*OWL 2 Specification*]) are satisfied in *I* w.r.t. *O*.

2.4 Models

Given a datatype map *D*, an interpretation $I = (\Delta_I , \Delta_D , \cdot^C , \cdot^{OP} , \cdot^{DP} , \cdot^I , \cdot^{DT} , \cdot^{LT} , \cdot^{FA})$ for *D* is a *model* of an OWL 2 ontology *O* w.r.t. *D* if an interpretation $J = (\Delta_I , \Delta_D , \cdot^C , \cdot^{OP} , \cdot^{DP} , \cdot^J , \cdot^{DT} , \cdot^{LT} , \cdot^{FA})$ for *D* exists such that $\cdot^J$ coincides with $\cdot^I$ on all named individuals and *J* satisfies *O*.

Thus, an interpretation *I* satisfying *O* is also a model of *O*. In contrast, a model *I* of *O* may not satisfy *O* directly; however, by modifying the interpretation of anonymous individuals, *I* can always be coerced into an interpretation *J* that satisfies *O*.

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2.5 Inference Problems

Let *D* be a datatype map and *V* a vocabulary over *D*. Furthermore, let *O* and O_1 be OWL 2 ontologies, *CE*, CE_1, and CE_2 class expressions, and *a* a named individual, such that all of them refer only to the vocabulary elements in *V*. Furthermore, *variables* are symbols that are not contained in *V*. Finally, a *Boolean conjunctive query* Q is a closed formula of the form

$$\exists\ x_1 , \ldots , x_n , y_1 , \ldots , y_m : [A_1 \wedge \ldots \wedge A_k]$$

where each A_i is an *atom* of the form `C(s)`, `OP(s,t)`, or `DP(s,u)` with `C` a class, `OP` an object property, `DP` a data property, `s` and `t` individuals or some variable x_j, and `u` a literal or some variable y_j.

The following inference problems are often considered in practice.

Ontology Consistency: *O* is *consistent* (or *satisfiable*) w.r.t. *D* if a model of *O* w.r.t. *D* and *V* exists.

Ontology Entailment: *O entails* O_1 w.r.t. *D* if every model of *O* w.r.t. *D* and *V* is also a model of O_1 w.r.t. *D* and *V*.

Ontology Equivalence: *O* and O_1 are *equivalent* w.r.t. *D* if *O* entails O_1 w.r.t. *D* and O_1 entails *O* w.r.t. *D*.

Ontology Equisatisfiability: *O* and O_1 are *equisatisfiable* w.r.t. *D* if *O* is satisfiable w.r.t. *D* if and only if O_1 is satisfiable w.r.t *D*.

Class Expression Satisfiability: *CE* is satisfiable w.r.t. *O* and *D* if a model $I = (\Delta_I , \Delta_D , \cdot^C , \cdot^{OP} , \cdot^{DP} , \cdot^I , \cdot^{DT} , \cdot^{LT} , \cdot^{FA})$ of *O* w.r.t. *D* and *V* exists such that $(CE)^C \neq \emptyset$.

Class Expression Subsumption: CE_1 is *subsumed* by a class expression CE_2 w.r.t. *O* and *D* if $(CE_1)^C \subseteq (CE_2)^C$ for each model $I = (\Delta_I , \Delta_D , \cdot^C , \cdot^{OP} , \cdot^{DP} , \cdot^I , \cdot^{DT} , \cdot^{LT} , \cdot^{FA})$ of *O* w.r.t. *D* and *V*.

Instance Checking: *a* is an *instance* of *CE* w.r.t. *O* and *D* if $(a)^I \in (CE)^C$ for each model $I = (\Delta_I , \Delta_D , \cdot^C , \cdot^{OP} , \cdot^{DP} , \cdot^I , \cdot^{DT} , \cdot^{LT} , \cdot^{FA})$ of *O* w.r.t. *D* and *V*.

Boolean Conjunctive Query Answering: *Q* is an *answer* w.r.t. *O* and *D* if *Q* is true in each model of *O* w.r.t. *D* and *V* according to the standard definitions of first-order logic.

In order to ensure that ontology entailment, class expression satisfiability, class expression subsumption, and instance checking are decidable, the following restriction w.r.t. *O* needs to be satisfied:

Each class expression of type **MinObjectCardinality**, **MaxObjectCardinality**, **ExactObjectCardinality**, and **ObjectHasSelf** that occurs in O_1, *CE*, CE_1, and CE_2 can contain only object property expressions that are simple in the axiom closure *Ax* of *O*.

For ontology equivalence to be decidable, O_1 needs to satisfy this restriction w.r.t. *O* and vice versa. These restrictions are analogous to the first condition from Section 11.2 of the OWL 2 Specification [*OWL 2 Specification*].

3 Independence of the Direct Semantics from the Datatype Map in OWL 2 DL (Informative)

OWL 2 DL has been defined so that the consequences of an OWL 2 DL ontology *O* do not depend on the choice of a datatype map, as long as the datatype map chosen contains all the datatypes occurring in *O*. This statement is made precise by the following theorem, and it has several useful consequences:

- One can apply the direct semantics to an OWL 2 DL ontology *O* by considering only the datatypes explicitly occurring in *O*.
- When referring to various reasoning problems, the datatype map *D* need not be given explicitly, as it is sufficient to consider an implicit datatype map containing only the datatypes from the given ontology.
- OWL 2 DL reasoners can provide datatypes not explicitly mentioned in this specification without fear that this will change the meaning of OWL 2 DL ontologies not using these datatypes.

Theorem DS1. Let O_1 and O_2 be OWL 2 DL ontologies over a vocabulary V and $D = (N_{DT} , N_{LS} , N_{FS} , \cdot^{DT} , \cdot^{LS} , \cdot^{FS})$ a datatype map such that each datatype mentioned in O_1 and O_2 is *rdfs:Literal*, a datatype defined in the respective ontology, or it occurs in N_{DT}. Furthermore, let $D' = (N_{DT}' , N_{LS}' , N_{FS}' , \cdot^{DT'} , \cdot^{LS'} , \cdot^{FS'})$ be a datatype map such that $N_{DT} \subseteq N_{DT}'$, $N_{LS}(DT) = N_{LS}'(DT)$, and $N_{FS}(DT) = N_{FS}'(DT)$ for each $DT \in N_{DT}$, and $\cdot^{DT'}$, $\cdot^{LS'}$, and $\cdot^{FS'}$ are extensions of $\cdot^{DT}$, $\cdot^{LS}$, and $\cdot^{FS}$, respectively. Then, O_1 entails O_2 w.r.t. D if and only if O_1 entails O_2 w.r.t. D'.

Proof. Without loss of generality, one can assume O_1 and O_2 to be in negation-normal form [*Description Logics*]. Furthermore, since datatype definitions in O_1 and O_2 are acyclic, one can assume that each defined datatype has been recursively replaced with its definition; thus, all datatypes in O_1 and O_2 are from $N_{DT} \cup \{$ *rdfs:Literal* $\}$. The claim of the theorem is equivalent to the following statement: an interpretation I w.r.t. D and V exists such that O_1 is and O_2 is not satisfied in I if and only if an interpretation I' w.r.t. D' and V exists such that O_1 is and O_2 is not satisfied in I'. The ($\Leftarrow$) direction is trivial since each interpretation I w.r.t. D' and V is also an interpretation w.r.t. D and V. For the ($\Rightarrow$) direction, assume that an interpretation $I = (\Delta_I , \Delta_D , \cdot^{C} , \cdot^{OP} , \cdot^{DP} , \cdot^{I} , \cdot^{DT} , \cdot^{LT} , \cdot^{FA})$ w.r.t. D and V exists such that O_1 is and O_2 is not satisfied in I. Let $I' = (\Delta_I , \Delta_D' , \cdot^{C'} , \cdot^{OP'} , \cdot^{DP'} , \cdot^{I'} , \cdot^{DT'} , \cdot^{LT'} , \cdot^{FA'})$ be an interpretation such that

- $\Delta_D{}'$ is obtained by extending Δ_D with the value space of all datatypes in $N_{DT}{}' \setminus N_{DT}$,
- $\cdot^{C\,'}$ coincides with $\cdot^{C}$ on all classes, and
- $\cdot^{DP\,'}$ coincides with $\cdot^{DP}$ on all data properties apart from *owl:topDataProperty*.

Clearly, $\mathtt{DataComplementOf(\ DR\ )}^{DT} \subseteq \mathtt{DataComplementOf(\ DR\ )}^{DT\,'}$ for each data range *DR* that is either a datatype, a datatype restriction, or an enumerated data range. The *owl:topDataProperty* property can occur in O_1 and O_2 only in tautologies. The interpretation of all other data properties is the same in *I* and *I'*, so $(CE)^{C} = (CE)^{C\,'}$ for each class expression *CE* occurring in O_1 and O_2. Therefore, O_1 is and O_2 is not satisfied in *I'*. QED

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